

# **Krawtchouk Polynomials Matrices and Transforms**

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Krawtchouk polynomials are formulated as matrices and properties of  
Krawtchouk transforms explored.

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# 1 Introduction

➔ **Krawtchouk polynomials** appear in a variety of contexts, most notably as orthogonal polynomials with respect to the binomial distribution.

- **Krawtchouk transform** on vectors.
- **Algorithm** for the Krawtchouk transform on vectors.
- **Krawtchouk expansions** of functions.
- **Operator calculus formulation** for the coefficients of Krawtchouk expansions.
- **Applications** of Krawtchouk transforms.

## 2 Krawtchouk Polynomials, Kravchuk Matrices

One may view Kravchuk matrices as an extension of the binomial coefficients. Consider the “degree-two algebraic rules” and translate them into a “second-degree Kravchuk matrix”:

$$\begin{aligned}(a + b)^2 &= a^2 + 2ab + b^2 \\(a + b)(a - b) &= a^2 - b^2 \\(a - b)^2 &= a^2 - 2ab + b^2\end{aligned}$$

read off

$$K^{(2)} = \begin{bmatrix} 1 & 1 & 1 \\ 2 & 0 & -2 \\ 1 & -1 & 1 \end{bmatrix}$$

The expansion coefficients make up the columns of the matrix.

## 2.1 Generating function

- The entries are determined by the expansion:

$$G(v; j, N) = (1 + v)^{N-j} (1 - v)^j = \sum_{i=0}^N v^i K_{ij}^{(N)}$$

- Expanding gives the explicit values of the matrix entries:

$$K_i(j, N) = K_{ij}^{(N)} = \sum_k (-1)^k \binom{j}{k} \binom{N-j}{i-k}$$

where matrix indices run from 0 to  $N$ .

► **Here** are the Kravchuk matrices of orders zero, one, and three:

$$K^{(0)} = \begin{bmatrix} 1 \end{bmatrix}$$

$$K^{(1)} = \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$K^{(3)} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 3 & 1 & -1 & -3 \\ 3 & -1 & -1 & 3 \\ 1 & -1 & 1 & -1 \end{bmatrix}$$

### 3 Interpretations

- **As** polynomials in  $j$  they are orthogonal with respect to the binomial distribution,  $\binom{N}{j}$ .
- **These** correspond to functionals of a random walk moving  $\pm 1$  with equal probabilities.
- **Transforms** of vectors correspond to expansions via matrices.
- **Transforms** of functions correspond to expansions in terms of polynomials.

## 4 Transform on Vectors

➔ Multiplying on the right by  $K$  gives the transform of  $\mathbf{f} = (f(0), f(1), \dots, f(N))$ . Multiply again by  $K$  using  $K^2 = 2^N \mathbf{I}$  to get the inverse transform.

$$\hat{\mathbf{f}} = \mathbf{f} K \quad \text{implies} \quad \mathbf{f} = 2^{-N} \hat{\mathbf{f}} K$$

• Explicitly, this is the expansion of the vector  $\mathbf{f}$  in terms of Krawtchouk polynomials in the variable  $j$ .

$$f(j) = 2^{-N} \sum_i \hat{f}(i) K_i(j, N)$$

• We have developed an algorithm for carrying out the transform.

## 4.1 Algorithm

➔ Given  $N > 0$ . Do the following for  $n = 0$  to  $N$  :

▶ **Step 0.** Given a row vector of length  $N + 1$ .

▶ **Step  $n$ .** You have  $n$  **current** rows.

Form  $n$  **new** rows by **summing** adjacent values.

Form the  $n + 1^{\text{st}}$  row by **differencing** adjacent values of the current  $n^{\text{th}}$  row.

● At step  $n$ , you have  $n + 1$  rows and  $N + 1 - n$  columns.

● After step  $N$ , you have a single column of  $N + 1$  values.

Transposed it is the Krawtchouk transform of the original row.

➔ Take the column that resulted from applying the algorithm as your new row. Apply the algorithm again.

Divide the result by  $2^N$  and you recover your original values.





## Examples

- Let  $N = 3$ . Start with 4, 2, 0, -3. Then we have

$$[4 \quad 2 \quad 0 \quad -3] \Rightarrow \begin{bmatrix} 6 & 2 & -3 \\ 2 & 2 & 3 \end{bmatrix} \Rightarrow \begin{bmatrix} 8 & -1 \\ 4 & 5 \\ 0 & -1 \end{bmatrix} \Rightarrow \begin{bmatrix} 7 \\ 9 \\ -1 \\ 1 \end{bmatrix}$$

- Start with a row of  $K^{(N)}$ , you get  $2^N$  times a vector with 1 in the corresponding spot.

$$[3 \quad 1 \quad -1 \quad -3] \Rightarrow 2^3 [0 \quad 1 \quad 0 \quad 0]$$

- Take a vector that starts with a binomial row,  $\binom{n}{i}$ .

Multiply on the left by  $K^{(N)}$ . It produces  $2^n$  times a binomial row with index  $N - n$ .

$$K^{(5)} [1 \quad 3 \quad 3 \quad 1 \quad 0 \quad 0]^t = 2^3 [1 \quad 2 \quad 1 \quad 0 \quad 0 \quad 0]^t$$

## 5 Expansions of Functions

In the random walk interpretation,  $j$  is the number of jumps to the left. The position  $x = N - 2j$ .

- The **generating function** for functions of  $x$  is

$$(1 + v)^{(N+x)/2} (1 - v)^{(N-x)/2} = \sum_{n \geq 0} \frac{v^n}{n!} K_n(x, N)$$

- A polynomial function of  $x$  of degree at most  $N$  has an expansion

$$f(x) = \sum_{0 \leq n \leq N} \tilde{f}(n) K_n(x, N)$$

- The coefficient  $\tilde{f}(n)$  has the operator calculus expression

$$\tilde{f}(n) = \frac{1}{n!} (\cosh D)^{N-n} (\sinh D)^n f(0)$$

where  $e^{\pm D} f(x) = f(x \pm 1)$ , shift operators on functions of  $x$ .

## 5.1 Operator calculus via Matrices

We can use the matrix of the operator  $D$  acting on the powers of  $x$  for symbolic calculation. Since  $D$  is nilpotent acting on polynomials in  $x$ , the exponentials reduce to finite sums. For example, take  $N = 4$ .

$$\hat{D} = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 2 & 0 & 0 \\ 0 & 0 & 0 & 3 & 0 \\ 0 & 0 & 0 & 0 & 4 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

We have  $\cosh \hat{D}$  and  $\tanh \hat{D}$  respectively:

$$\begin{bmatrix} 1 & 0 & 1 & 0 & 1 \\ 0 & 1 & 0 & 3 & 0 \\ 0 & 0 & 1 & 0 & 6 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}, \begin{bmatrix} 0 & 1 & 0 & -2 & 0 \\ 0 & 0 & 2 & 0 & -8 \\ 0 & 0 & 0 & 3 & 0 \\ 0 & 0 & 0 & 0 & 4 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$



## Example

For  $N=4$ , let  $f(x) = x^4 + 2x^3 - x^2 + 5x$ .

We find, with  $N = 4$ , that

$$f = K_4 + 2K_3 + 15K_2 + 25K_1 + 36$$

where  $K_0 = 1$ ,

$$K_1 = x,$$

$$K_3 = x^3 - 10x$$

$$K_2 = x^2 - 4,$$

$$K_4 = x^4 - 16x^2 + 24$$

This is obtained by multiplying the column of coefficients of  $f(x)$  by the matrix  $Y$  formed by the top rows of  $(\cosh^N \hat{D})(\tanh^n \hat{D})/n!$ , for  $0 \leq n \leq N$ , which are readily computed iteratively. In this example we have

$$Y = \begin{bmatrix} 1 & 0 & 4 & 0 & 40 \\ 0 & 1 & 0 & 10 & 0 \\ 0 & 0 & 1 & 0 & 16 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

## 6 Further aspects

- **General**  $p, q$ . Polynomials with parameters  $p$  and  $q$  arise from Bernoulli trials where the probability of success is  $p$ , with  $q = 1 - p$ . They arise as well when working over finite fields, in which case  $q$  is the number of elements of the field.
- **Multivariate** polynomials are orthogonal with respect to corresponding multinomial distributions. Functions of several variables correspond to random walks in higher dimensions.
- **Positivity** results hold for transforms of polynomial functions.
- **Variety** of applications is seen in the references.

## References

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